# Extensions of embeddings in the $\Sigma_{2}^{0}$ enumeration degrees 

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## Enumeration reducibility

Definition (various authors, 1950s)
For (nonempty) $A, B \subseteq \mathbb{N}$, we say that $A$ is enumeration reducible to $B\left(A \leq_{\mathrm{e}} B\right)$ if
every enumeration of $B$ computes an enumeration of $A$.

Example
Let $B$ be a maximal independent set of vertices in a computable graph. Then $B^{c} \leq_{\mathrm{e}} B$ (but $B \not \leq_{\mathrm{e}} B^{c}$ in general.)

The enumeration degrees (e-degrees) are defined from $\leq_{\mathrm{e}}$ in the same way that the Turing degrees are defined from $\leq_{T}$.

They form an upper-semilattice under $\leq_{e}$ and the usual effective join.

## Beyond the Turing degrees: The enumeration degrees

The Turing degrees embed into the e-degrees in a natural way:

$$
A \leq_{\mathrm{T}} B \quad \text { if and only if } A \oplus A^{c} \leq_{\mathrm{e}} B \oplus B^{c}
$$

so the map $A \mapsto A \oplus A^{c}$ induces an embedding.

## Questions we can ask about a partial order

1. Is it linear?
2. Which finite partial orders embed into it?
3. Is it dense?
4. Given finite partial orders $\mathcal{P} \subseteq \mathcal{Q}$, can every embedding of $\mathcal{P}$ into it be extended to an embedding of $\mathcal{Q}$ ?
5. (!) Given a first-order sentence in the language of $\{\leq\}$, can we algorithmically decide if it is true?

For many degree structures, the answer to 5 is very much no (Slaman, Woodin 1997).

## A central question

On the other hand, every finite partial order embeds into the e-degrees (corollary of Sacks 1963), so we can compute if a sentence of the form

$$
\exists a_{0} \exists a_{1} \cdots \exists a_{n} \text { (Boolean combination of } a_{i} \leq a_{j} \text { ) }
$$

holds in the e-degrees by checking whether the Boolean combination is consistent with the axioms of partial orders.

At what point does computability break down?

## A countable substructure: The $\Sigma_{2}^{0}$ e-degrees

We are working on this question for the $\Sigma_{2}^{0}$ e-degrees.

Reasons to study the $\Sigma_{2}^{0}$ e-degrees:

1. They are analogous to the c.e. Turing degrees
2. Any nontrivial automorphism of the e-degrees must move some $\Sigma_{2}^{0}$ e-degree (Slaman, Soskova 2017)
3. They exhibit unusual order-theoretic phenomenon (next slide)

## The $\Sigma_{2}^{0}$ e-degrees: Ahmad pairs

The $\Sigma_{2}^{0}$ e-degrees are dense (Cooper).
Compare: The c.e. Turing degrees are dense (Sacks density).

## Question (Cooper)

Does the $\Sigma_{2}^{0}$ e-degrees satisfy the same first-order sentences as the c.e. Turing degrees?

Theorem (Ahmad 1989) In the $\Sigma_{2}^{0}$ e-degrees, there are incomparable $\mathbf{a}$ and $\mathbf{b}$ such that if $\mathbf{x}<_{e} \mathbf{a}$, then $\mathbf{x} \leq_{e} \mathbf{b}$.


Notice that a cannot be the join of two degrees below it.
On the other hand, in the c.e. Turing degrees, every degree is the join of two degrees below it (Sacks splitting).

## There are no Ahmad triples

Theorem (Ahmad 1989)
There are incomparable $\mathbf{a}$ and $\mathbf{b}$ such that if $\mathbf{x}<_{\mathrm{e}} \mathbf{a}$, then $\mathbf{x} \leq_{\mathrm{e}} \mathbf{b}$.

Theorem (G., Lempp, Ng,
Soskova 2021)
For every incomparable $\mathbf{a}, \mathbf{b}, \mathbf{c}$, there is some $\mathbf{x}$ such that:

- $\mathbf{x}<_{\mathrm{e}} \mathbf{a}$ but $\mathbf{x} \not Z_{\mathrm{e}} \mathbf{b}$, OR
- $\mathbf{x}<_{\mathrm{e}} \mathbf{b}$ but $\mathbf{x} \not \mathbb{e}_{\mathrm{e}} \mathbf{c}$.
(Actually we just need $\mathbf{a} \not \leq \mathbf{b}$ and $\mathbf{b} \not \leq \mathbf{c}$, so a and could be comparable or even equal.)



## Reformulating as extensions of embeddings

Theorem (Ahmad 1989)
There are incomparable $\mathbf{a}$ and $\mathbf{b}$ such that if $\mathbf{x}<_{e} \mathbf{a}$, then $\mathbf{x} \leq_{e} \mathbf{b}$.
Reformulation:

> Not every embedding of the antichain $\{a, b\}$ (into the $\Sigma_{2}^{0}$ e-degrees) can be extended to an embedding of $\{a, b, x\}$ where $x<a$ and $x \not \leq b$.

Our result can be reformulated similarly:
Every embedding of the antichain $\{a, b, c\}$ can be extended to one of the following:

- an embedding of $\{a, b, c, x\}$ where $x<a$ and $x \not \leq b$
- an embedding of $\{a, b, c, x\}$ where $x<b$ and $x \not \leq c$.
(Actually, there are four choices here because in the first ordering, we didn't specify the relationship between $x$ and
$c$, and in the second ordering, we didn't specify the relationship between $x$ and a.)


## "Disjunctive" results

Our result on no "Ahmad triples" generalizes the following

Theorem (Ahmad 1989)
There are no "symmetric Ahmad pairs", i.e., for every incomparable $\mathbf{a}$ and $\mathbf{b}$, there is some $\mathbf{x}$ such that:

- $\mathbf{x}<_{\mathrm{e}} \mathbf{a}$ but $\mathbf{x} \not \mathbb{Z}_{\mathrm{e}} \mathbf{b}$, or
$-\mathbf{x}<_{\mathrm{e}} \mathbf{b}$ but $\mathbf{x} \not \mathbb{Z}_{\mathrm{e}} \mathbf{a}$.

In other words, every embedding of the antichain $\{a, b\}$ can be extended to one of the following:

- an embedding of $\{a, b, x\}$ where $x<a$ and $x \not \leq b$
- an embedding of $\{a, b, x\}$ where $x<b$ and $x \not \leq a$.


## What's the goal here?

We'd like to give an algorithm for deciding the two-quantifier first-order theory of the $\Sigma_{2}^{0}$ e-degrees. In terms of quantifier alternations, two quantifiers is the most we can hope to decide (Kent 2006).

Fact: Every two-quantifier sentence can be thought of as a disjunctive extension of embeddings problem.

Our focus (for now) is the following special case:
One-point extensions of antichains problem
Given a (finite) antichain $\mathcal{P}$ and one-point extensions $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$ of $\mathcal{P}$, decide whether every embedding of $\mathcal{P}$ into the $\Sigma_{2}^{0}$ e-degrees extends to an embedding of some $\mathcal{Q}_{i}$.
$n=0$ was solved by Lempp, Slaman, Sorbi 2005. $n \geq 1$ is open.

## One-point extensions of antichains problem

We'll restrict ourselves further to one-point extensions where the new element is not above any of the old elements.

Given an antichain $\mathcal{P}=\left\{a_{0}, \ldots, a_{k}\right\}$ and such one-point extensions $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$ of $\mathcal{P}$, let $x_{i}$ denote the new element added by $\mathcal{Q}_{i}$.

Suppose that in $\mathcal{Q}_{i}$, we have $x_{i}<a_{i}$ and $x_{i} \mid a_{j}$ for $j \neq i$.


We shall call such $\mathcal{Q}_{i}$ singleton extensions.

## Singleton extensions and weak Ahmad bases

If $\mathcal{Q}_{i}$ is a singleton extension and $\mathbf{a}_{0}, \ldots, \mathbf{a}_{k}$ does not extend to an embedding of $\mathcal{Q}_{i}$ then:

Every $\mathbf{x}<_{\mathrm{e}} \mathbf{a}_{i}$ must be below some other $\mathbf{a}_{j}$.
If $\mathbf{a}_{i}$ has the above property we say that it is a weak Ahmad base.

- If $\mathbf{a}$ and $\mathbf{b}$ form an Ahmad pair, then $\mathbf{a}$ is a weak Ahmad base.
- Not all weak Ahmad bases come from Ahmad pairs (G., Lempp, Ng, Soskova 2021).

Right: Every $\mathbf{x}<_{\mathrm{e}} \mathbf{a}_{0}$ is below $\mathbf{a}_{1}$ or below $\mathbf{a}_{2}$, but $\mathbf{a}_{0}$ does not form an Ahmad pair with $\mathbf{a}_{1}$ or with $\mathbf{a}_{2}$, as witnessed by $\mathbf{x}_{1}$ and $\mathbf{x}_{2}$ respectively.


We want to understand which degrees can be weak Ahmad bases, because in general, we may want to construct an embedding of $\mathcal{P}$ which cannot be extended to any of $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$, not just a single $\mathcal{Q}_{i}$.

Theorem (G., Lempp, Ng, Soskova ongoing)
If $\mathbf{a}$ and $\mathbf{b}$ form an Ahmad pair, then $\mathbf{b}$ is not a weak Ahmad base.
This generalizes our result on "no Ahmad triple".
In fact, we believe we have proved the following stronger statement (*):

If $\mathbf{a}_{0}, \ldots, \mathbf{a}_{\ell}$ are incomparable weak Ahmad bases, then
for each $i \leq \ell$, there is some $\mathbf{x}_{i}<{ }_{\mathrm{e}} \mathbf{a}_{i}$ such that
$\left.\mathbf{x}_{i}\right|_{\mathrm{e}} \mathbf{a}_{j}$ for all $j \neq i$.
(To derive the theorem from $(*)$, take $\ell=1, \mathbf{a}_{0}=\mathbf{a}, \mathbf{a}_{1}=\mathbf{b}$.)
(*): If $\mathbf{a}_{0}, \ldots, \mathbf{a}_{\ell}$ are incomparable weak Ahmad bases, then for each $i \leq \ell$, there is some $\mathbf{x}_{i}<_{\mathrm{e}} \mathbf{a}_{i}$ such that $\left.\mathbf{x}_{i}\right|_{\mathrm{e}} \mathbf{a}_{j}$ for all $j \neq i$.

Assuming $(*)$, we can derive a combinatorial condition on $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$ which is necessary for there to be an embedding of $\left\{a_{0}, \ldots, a_{k}\right\}$ which does not extend to any of $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$.


We must have that $\mathbf{x}_{0}, \mathbf{x}_{1}, \ldots, \mathbf{x}_{\ell}$ cannot be used to extend the embedding $\mathbf{a}_{0}, \ldots, \mathbf{a}_{k}$ to any of $\mathcal{Q}_{\ell+1}, \ldots, \mathcal{Q}_{n}$.

Our plan is to prove that this combinatorial condition is also sufficient for constructing an embedding of $\left\{a_{0}, \ldots, a_{k}\right\}$ which does not extend to any of $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$.

If true, this would give us an algorithm for deciding whether every embedding of a given finite antichain $\mathcal{P}$ extends to an embedding of some $\mathcal{Q}_{i}$, where $\mathcal{Q}_{0}, \ldots, \mathcal{Q}_{n}$ are one-point extensions of $\mathcal{P}$ where the new element is not above any old elements.

## Thanks!

