# Extensions of embeddings in the $\Sigma^0_2$ enumeration degrees

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# Enumeration reducibility

Definition (various authors, 1950s) For (nonempty)  $A, B \subseteq \mathbb{N}$ , we say that A is enumeration reducible to B ( $A \leq_{e} B$ ) if

every enumeration of B computes an enumeration of A.

#### Example

Let B be a maximal independent set of vertices in a computable graph. Then  $B^c \leq_{e} B$  (but  $B \not\leq_{e} B^c$  in general.)

The enumeration degrees (e-degrees) are defined from  $\leq_e$  in the same way that the Turing degrees are defined from  $\leq_T$ .

They form an upper-semilattice under  $\leq_{\rm e}$  and the usual effective join.

Beyond the Turing degrees: The enumeration degrees

The Turing degrees embed into the e-degrees in a natural way:  $A \leq_{\mathrm{T}} B$  if and only if  $A \oplus A^c \leq_{\mathrm{e}} B \oplus B^c$ so the map  $A \mapsto A \oplus A^c$  induces an embedding. Questions we can ask about a partial order

- 1. Is it linear?
- 2. Which finite partial orders embed into it?
- 3. Is it dense?
- Given finite partial orders P ⊆ Q, can every embedding of P into it be extended to an embedding of Q?
- 5. (!) Given a first-order sentence in the language of  $\{\leq\}$ , can we algorithmically decide if it is true?

For many degree structures, the answer to 5 is very much no (Slaman, Woodin 1997).

### A central question

On the other hand, every finite partial order embeds into the e-degrees (corollary of Sacks 1963), so we can compute if a sentence of the form

 $\exists a_0 \exists a_1 \cdots \exists a_n \text{ (Boolean combination of } a_i \leq a_j \text{)}$ 

holds in the e-degrees by checking whether the Boolean combination is consistent with the axioms of partial orders.

At what point does computability break down?

A countable substructure: The  $\Sigma_2^0$  e-degrees

We are working on this question for the  $\Sigma_2^0$  e-degrees.

Reasons to study the  $\Sigma_2^0$  e-degrees:

1. They are analogous to the c.e. Turing degrees

- 2. Any nontrivial automorphism of the e-degrees must move some  $\Sigma_2^0$  e-degree (Slaman, Soskova 2017)
- 3. They exhibit unusual order-theoretic phenomenon (next slide)

# The $\Sigma_2^0$ e-degrees: Ahmad pairs

The  $\Sigma_2^0$  e-degrees are dense (Cooper).

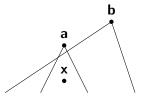
Compare: The c.e. Turing degrees are dense (Sacks density).

### Question (Cooper)

Does the  $\Sigma_2^0$  e-degrees satisfy the same first-order sentences as the c.e. Turing degrees?

### Theorem (Ahmad 1989)

In the  $\Sigma_2^0$  e-degrees, there are incomparable a and b such that if  $x <_{\rm e} a$ , then  $x \leq_{\rm e} b$ .



Notice that a cannot be the join of two degrees below it.

On the other hand, in the c.e. Turing degrees, every degree is the join of two degrees below it (Sacks splitting).

### There are no Ahmad triples

### Theorem (Ahmad 1989)

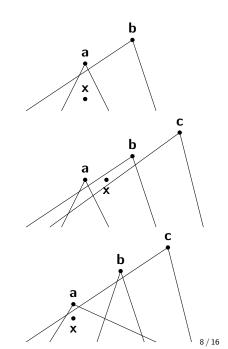
There are incomparable **a** and **b** such that if  $\mathbf{x} <_e \mathbf{a}$ , then  $\mathbf{x} \leq_e \mathbf{b}$ .

# Theorem (G., Lempp, Ng, Soskova 2021)

For every incomparable **a**, **b**, **c**, there is some **x** such that:

- ▶  $\mathbf{x} <_{e} \mathbf{a}$  but  $\mathbf{x} \not\leq_{e} \mathbf{b}$ , OR
- ▶  $\mathbf{x} <_{e} \mathbf{b}$  but  $\mathbf{x} \not\leq_{e} \mathbf{c}$ .

(Actually we just need  $\mathbf{a} \leq \mathbf{b}$  and  $\mathbf{b} \leq \mathbf{c}$ , so  $\mathbf{a}$  and  $\mathbf{c}$  could be comparable or even equal.)



# Reformulating as extensions of embeddings

### Theorem (Ahmad 1989)

There are incomparable **a** and **b** such that if  $\mathbf{x} <_{e} \mathbf{a}$ , then  $\mathbf{x} \leq_{e} \mathbf{b}$ .

Reformulation:

Not every embedding of the antichain  $\{a, b\}$  (into the  $\Sigma_2^0$  e-degrees) can be extended to an embedding of  $\{a, b, x\}$  where x < a and  $x \not\leq b$ .

Our result can be reformulated similarly:

Every embedding of the antichain  $\{a, b, c\}$  can be extended to one of the following:

an embedding of {a, b, c, x} where x < a and x ≤ b</li>
 an embedding of {a, b, c, x} where x < b and x ≤ c.</li>

(Actually, there are four choices here because in the first ordering, we didn't specify the relationship between x and c, and in the second ordering, we didn't specify the relationship between x and a.)

### "Disjunctive" results

Our result on no "Ahmad triples" generalizes the following

### Theorem (Ahmad 1989)

There are no "symmetric Ahmad pairs", i.e., for every incomparable  $\mathbf{a}$  and  $\mathbf{b}$ , there is some  $\mathbf{x}$  such that:

▶ x 
$${<_{\mathrm{e}}}$$
 a but x  ${\not\leq_{\mathrm{e}}}$  b, or

► x <
$$_{
m e}$$
 b but x  $\not\leq_{
m e}$  a.

In other words, every embedding of the antichain  $\{a, b\}$  can be extended to one of the following:

- ▶ an embedding of  $\{a, b, x\}$  where x < a and  $x \not\leq b$
- ▶ an embedding of  $\{a, b, x\}$  where x < b and  $x \not\leq a$ .

# What's the goal here?

We'd like to give an algorithm for deciding the **two-quantifier** first-order theory of the  $\Sigma_2^0$  e-degrees. In terms of quantifier alternations, two quantifiers is the most we can hope to decide (Kent 2006).

Fact: Every two-quantifier sentence can be thought of as a disjunctive extension of embeddings problem.

Our focus (for now) is the following special case:

One-point extensions of antichains problem Given a (finite) **antichain**  $\mathcal{P}$  and **one-point** extensions  $\mathcal{Q}_0, \ldots, \mathcal{Q}_n$ of  $\mathcal{P}$ , decide whether every embedding of  $\mathcal{P}$  into the  $\Sigma_2^0$  e-degrees extends to an embedding of some  $\mathcal{Q}_i$ .

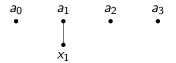
n = 0 was solved by Lempp, Slaman, Sorbi 2005.  $n \ge 1$  is open.

### One-point extensions of antichains problem

We'll restrict ourselves further to one-point extensions where the new element is **not above** any of the old elements.

Given an antichain  $\mathcal{P} = \{a_0, \ldots, a_k\}$  and such one-point extensions  $\mathcal{Q}_0, \ldots, \mathcal{Q}_n$  of  $\mathcal{P}$ , let  $x_i$  denote the new element added by  $\mathcal{Q}_i$ .

Suppose that in  $Q_i$ , we have  $x_i < a_i$  and  $x_i \mid a_j$  for  $j \neq i$ .



We shall call such  $Q_i$  singleton extensions.

### Singleton extensions and weak Ahmad bases

If  $Q_i$  is a singleton extension and  $\mathbf{a}_0, \ldots, \mathbf{a}_k$  does not extend to an embedding of  $Q_i$  then:

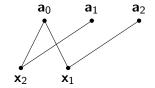
Every  $\mathbf{x} <_{e} \mathbf{a}_{i}$  must be below some other  $\mathbf{a}_{i}$ .

If  $\mathbf{a}_i$  has the above property we say that it is a weak Ahmad base.

▶ If **a** and **b** form an Ahmad pair, then **a** is a weak Ahmad base.

 Not all weak Ahmad bases come from Ahmad pairs (G., Lempp, Ng, Soskova 2021).

Right: Every  $\mathbf{x} <_{e} \mathbf{a}_{0}$  is below  $\mathbf{a}_{1}$  or below  $\mathbf{a}_{2}$ , but  $\mathbf{a}_{0}$  does not form an Ahmad pair with  $\mathbf{a}_{1}$  or with  $\mathbf{a}_{2}$ , as witnessed by  $\mathbf{x}_{1}$  and  $\mathbf{x}_{2}$  respectively.



We want to understand which degrees can be weak Ahmad bases, because in general, we may want to construct an embedding of  $\mathcal{P}$  which cannot be extended to any of  $\mathcal{Q}_0, \ldots, \mathcal{Q}_n$ , not just a single  $\mathcal{Q}_i$ .

### Theorem (G., Lempp, Ng, Soskova ongoing)

If  $\mathbf{a}$  and  $\mathbf{b}$  form an Ahmad pair, then  $\mathbf{b}$  is not a weak Ahmad base.

This generalizes our result on "no Ahmad triple".

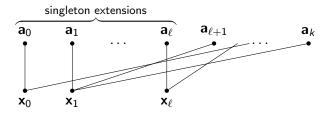
In fact, we believe we have proved the following stronger statement (\*):

If  $\mathbf{a}_0, \ldots, \mathbf{a}_\ell$  are incomparable weak Ahmad bases, then for each  $i \leq \ell$ , there is some  $\mathbf{x}_i <_{\mathrm{e}} \mathbf{a}_i$  such that  $\mathbf{x}_i \mid_{\mathrm{e}} \mathbf{a}_j$  for all  $j \neq i$ .

(To derive the theorem from (\*), take  $\ell = 1$ ,  $\mathbf{a}_0 = \mathbf{a}$ ,  $\mathbf{a}_1 = \mathbf{b}$ .)

(\*): If  $\mathbf{a}_0, \ldots, \mathbf{a}_\ell$  are incomparable weak Ahmad bases, then for each  $i \leq \ell$ , there is some  $\mathbf{x}_i <_{\mathrm{e}} \mathbf{a}_i$  such that  $\mathbf{x}_i \mid_{\mathrm{e}} \mathbf{a}_j$  for all  $j \neq i$ .

Assuming (\*), we can derive a combinatorial condition on  $Q_0, \ldots, Q_n$  which is **necessary** for there to be an embedding of  $\{a_0, \ldots, a_k\}$  which does not extend to any of  $Q_0, \ldots, Q_n$ .



We must have that  $\mathbf{x}_0, \mathbf{x}_1, \dots, \mathbf{x}_{\ell}$  cannot be used to extend the embedding  $\mathbf{a}_0, \dots, \mathbf{a}_k$  to any of  $\mathcal{Q}_{\ell+1}, \dots, \mathcal{Q}_n$ .

Our plan is to prove that this combinatorial condition is also **sufficient** for constructing an embedding of  $\{a_0, \ldots, a_k\}$  which does not extend to any of  $Q_0, \ldots, Q_n$ .

If true, this would give us an algorithm for deciding whether every embedding of a given finite antichain  $\mathcal{P}$  extends to an embedding of some  $\mathcal{Q}_i$ , where  $\mathcal{Q}_0, \ldots, \mathcal{Q}_n$  are one-point extensions of  $\mathcal{P}$ where the new element is not above any old elements.

# Thanks!